

# Constant Propagation

A More Complex Semilattice  
A Nondistributive Framework

# The Point

- ◆ Instead of doing constant folding by RD's, we can maintain information about what constant, if any, a variable has at each point.
- ◆ An interesting example of a DF framework not of the gen-kill type.
- ◆ A simple version of static type analysis.

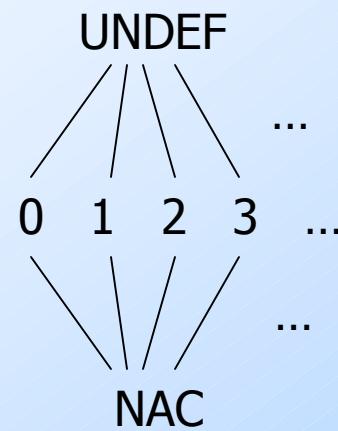
# Domain of Values

- ◆ The set of values propagated is the set of mappings from variables to values of their type.
- ◆ Example: [x → 5, s → "cat", y → UNDEF, z → NAC]
  - ◆ UNDEF = "We don't yet know anything."
  - ◆ NAC = "Not a constant" = we know too much for any constant to satisfy."

# The Semilattice

- ◆ A *product lattice*, one component for each variable.
- ◆ Each component lattice consists of:
  1. UNDEF (the top element).
  2. NAC (the bottom element).
  3. All values from a type, e.g., integers, strings.

# Picture



# The Meet Operation

- ◆ The diagram represents  $\leq$ . That is:
  1. Any constant  $\leq$  UNDEF.
  2. NAC  $\leq$  any constant.
- ◆ Equivalently, for any constants x and y:
  1. UNDEF  $\wedge$  x = x.
  2. NAC  $\wedge$  x = NAC.
  3. NAC  $\wedge$  UNDEF = NAC.
  4.  $x \wedge x = x$  but  $x \wedge y = \text{NAC}$  if  $x \neq y$ .

# The Product Lattice

- ◆ Call each of the lattices just described a *diamond lattice*.
- ◆ The lattices we use are products of diamond lattices.
- ◆ For the product  $D_1 * D_2 * \dots * D_n$ , the values are  $[v_1, v_2, \dots, v_n]$ , where each  $v_i$  is in  $D_i$ .

# Meet in Product Lattices

- ◆  $[v_1, v_2, \dots, v_n] \wedge [w_1, w_2, \dots, w_n] = [v_1 \wedge w_1, v_2 \wedge w_2, \dots, v_n \wedge w_n] =$   
*componentwise meet.*
- ◆ In terms of  $\leq$ :  
 $[v_1, v_2, \dots, v_n] \leq [w_1, w_2, \dots, w_n]$   
if and only if  $v_i \leq w_i$  for all  $i$ .

# Intuitive Meaning

1. If variable  $x$  is mapped to UNDEF (i.e., in the product-lattice value, the component for  $x$  is UNDEF), then we do not know anything about  $x$ .
2. If  $x$  is mapped to constant  $c$ , then we only know of paths where  $x$  has value  $c$ .
3. If  $x$  is mapped to NAC, we know about paths where  $x$  has different values.

# Product-Lattice Values as Mappings

- ◆ Think of a lattice element as a mapping from variables to values {UNDEF, NAC, constants}.
- ◆ Lattice element is  $m$ , and  $m(x)$  is the value to which  $m$  maps variable  $x$ .

# Transfer Functions --- (1)

- ◆ Transfer functions map lattice elements to lattice elements.
- ◆ Suppose  $m$  is the variable->constant mapping just before a statement  
 $x = y + z.$
- ◆ Let  $f(m) = m'$  be the transfer function associated with  $x = y + z.$

# Transfer Functions --- (2)

- ◆ If  $m(y) = c$  and  $m(z) = d$ , then  $m'(x) = c+d$ .
- ◆ If  $m(y) = \text{NAC}$  or  $m(z) = \text{NAC}$ , then  $m'(x) = \text{NAC}$ .
- ◆ Otherwise, if  $m(y) = \text{UNDEF}$  or  $m(z) = \text{UNDEF}$ , then  $m'(x) = \text{UNDEF}$ .
- ◆  $m'(w) = m(w)$  for all  $w$  other than  $x$ .

# Transfer Functions --- (3)

- ◆ Similar rules for other types of statements (see text).
- ◆ For a block, compose the transfer functions of the individual statements.

# Iterative Algorithm

- ◆ It's a plain-ol' Forward iteration, with the meet and transfer functions as given.
- ◆ The framework is monotone and has bounded depth, so it converges to a safe solution.

# Finite Depth

- ◆ The value of any IN or OUT can only decrease.
  - ◆ Verify from transfer functions (monotonicity).
- ◆ Values are finite-length vectors, and each component can only decrease twice.
  - ◆ From UNDEF to a constant to NAC.
- ◆ If no IN or OUT decreases in any component in a round, we stop.

# Monotonicity --- (1)

- ◆ Need to show  $m \leq n$  implies  $f(m) \leq f(n)$ .
- ◆ Show for function  $f$  associated with a single statement.
- ◆ Composition of monotone functions is monotone.
- ◆ That's enough to show monotonicity for all possible transfer functions.

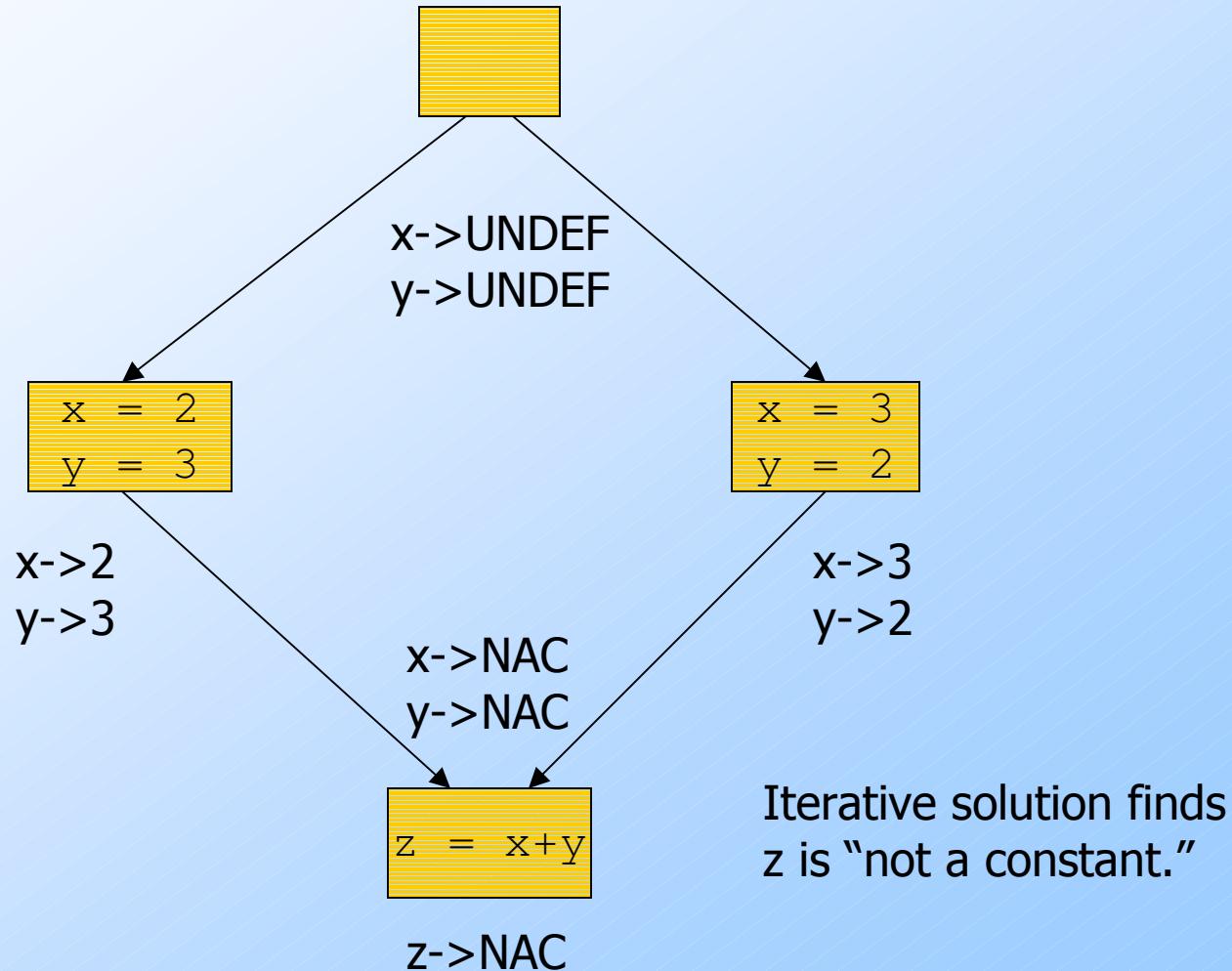
# Monotonicity --- (2)

- ◆ **One case:** let  $f$  be the function associated with  $x = y+z$ .
- ◆ **One subcase:**  $m(y) = c; m(z) = d; n(y) = c; n(z) = \text{UNDEF}; m(w) = n(w)$  otherwise. Thus,  $m \leq n$ .
- ◆ Then  $(f(m))(x) = c+d$  and  $(f(n))(x) = \text{UNDEF}$ .
- ◆ Thus  $(f(m))(w) \leq (f(n))(w)$  for all  $w$ .

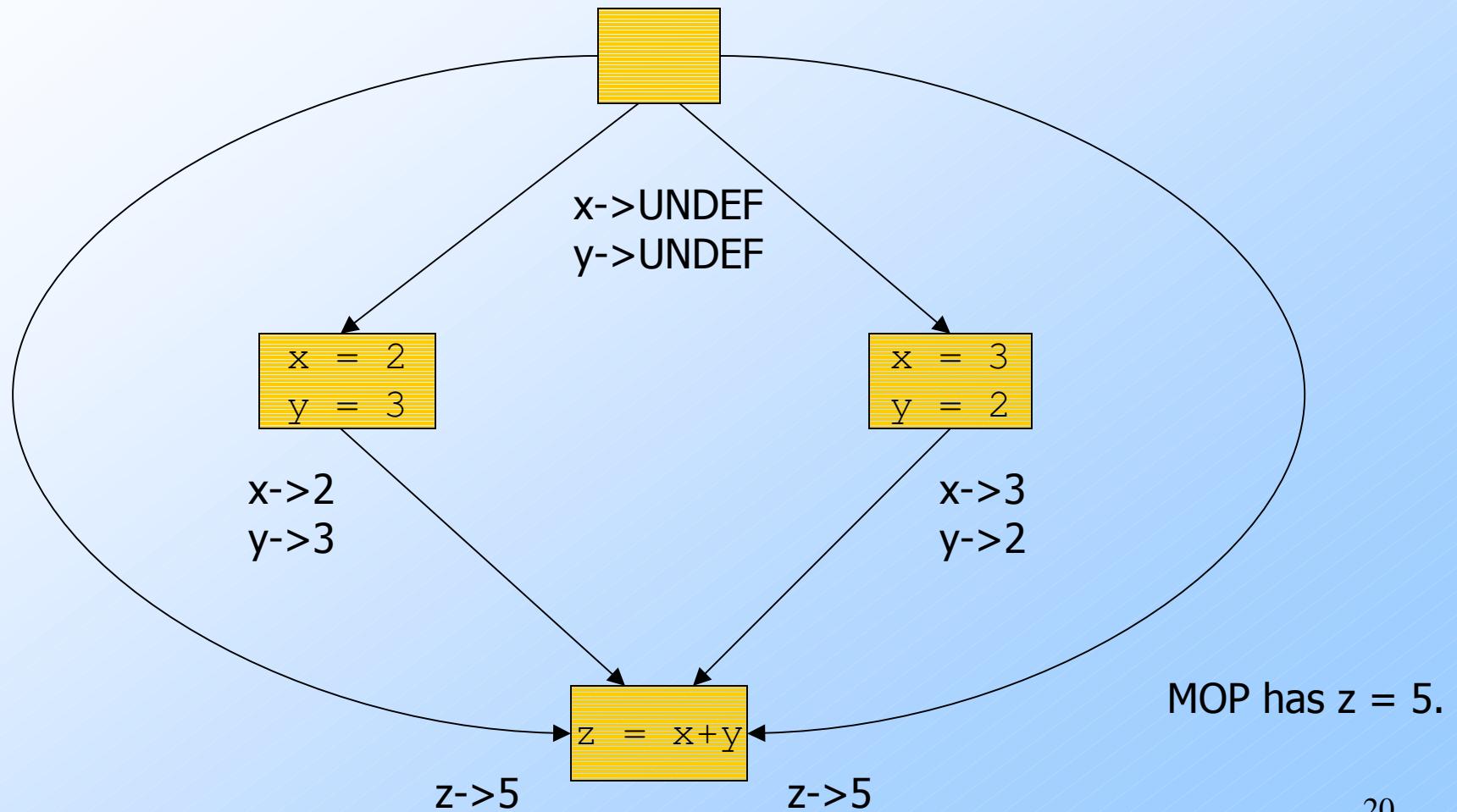
# Nondistributivity

- ◆ First example of a framework that is not distributive.
- ◆ Thus, iterative solution is not the MOP.
- ◆ We'll show an example where MFP appears to include impossible paths.

# Example: Nondistributivity



## Example: Nondistributivity --- (2)



# Example: Nondistributivity --- (3)

- ◆ We observe that MFP differs from the MOP solution.
- ◆ That proves the framework is not distributive.
  - ◆ Because every distributive framework has  $MFP = MOP$ .