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NOTES ON A PROBLEM INVOLVING PERMUTATIONS AS SUBSEQUENCES

BY

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ABSTRACT :

The problem (attributed to R. M. Karp by Knuth (see #36 of [1])) is to describe the sequences of minimum length which contain, as subsequences, all the permutations of an alphabet of n symbols. This paper catalogs some of the easy observations on the problem and proves that the minimum lengths for n-5, n=6 6 n=7 are 19, 28 and 39 respectively.' Also presented is a construction which yields (for n>2) many appropriate sequences of length n^2-2n+4 so giving an upper bound on length of minimum strings which matches exactly all known values,

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1 NOTATION.

a) Let S be a sequence of symbols. $|S|$ will be used to denote the total number of symbols in S and so we observe, for example, $|x\ y\ x\ z| = 4$.

b) We say $x \subset y$ in the case where x is a subsequence of y and we say "x is equivalent to y" if x can be obtained from y by a simple change of alphabet; we denote this equivalence by '$\equiv$'.
(e.g. $x\ y \subset x\ y\ y\ x$, $x\ y\ z\ x \equiv 1\ 2\ 3\ 1$)

c) $P(A)$ is used to denote the set of sequences uhich are permutations of an alphabet A. Cardinality of $P(A)$ will be $(|A|)!$. Also, $P'(A,n)$ is is the set of permutations of all sub-alphabets of A of size n (where $n \le |A|$). Clearly, $P(A)=P'(A,|A|)$.

d) If A is an alphabet then $Q(A)=\{x \mid x \in A' \wedge \forall y.(y \in P(A) \supset y \subset x)\}$ where A' is the set of sequences over alphabet A. For example, abcacba $\in Q(abc)$. Also, $Q'(A,n)$ is taken to be the set $\{x \mid x \in A' \wedge \forall y.(y \in P(A,n) \supset y \subset x)\}$. So, for example, zyxwxyz $\in Q'(wxyz,2)$.

e) Now, the LENGTHS of the shortest sequences in $Q(A)$ and $Q'(A,n)$ depend only on the SIZE of the alphabet A. Hence, take $M(n)$ to be the length of the shortest sequence in $Q(1\ 2\ 3 \ldots n)$ and $M'(n,m)$ to be the length of the shortest sequence in $Q'(1\ 2\ 3\ \ldots\ n, m)$.
So, for example, $M(1)=1$, $M(2)=3$ and $M'(n,1)=n$.

f) $S(n)$ denotes the n-th symbol of sequence S.
$S(n:m)$ denotes that contiguous subsequence of sequence S which is the symbols from position number n in S to position number m.
$\#(S,x)$ denotes the number of ocurrences of the symbol x in sequence S.

g) "CPAF X" is just an abbreviation for "Consider the Permutations of the current Alphabet of the Form X". The greek letters uhich appear in X denote arbitrary sequences of symbols.
For example, if the alphabet under discussion were abcde, the command "CPAF bαc" would mean "Consider Permutations of abcde which start with b and end with c".

2 SOME EASY OBSERVATIONS.

2.1 $M(1)=1$.

2.2 $M(2)=3$.

2.3 $M(3)=7$.

2.4 $M'(n,1)=n$.

2.5 $M'(n,2)=(2n-1)$ can be seen as follows:

$M'(n,2)\le 2n-1$ since if A is an alphabet of length n, then the sequence $AA(2:2n)$ is a member of $Q'(A,2)$.
$M'(n,2)\ge 2n-1$ since if A is an alphabet of size n, S is a member of $Q'(A,2)$ and $|S|<2n-1$ then at least two of the symbols of A (x and y, say) only appear once in S; hence 1 of the sequences 'xy' and 'yx' are not subsequences of S.

2.6 $M'(n,m)\ge(m.(2n-m+1)/2)$ ($n\ge m$, of course)

This result is more easily remembered as

$$M'(n,m)\ge n + n\text{-}l \text{ t } n\text{-}2 \text{ t } \ldots \text{ t } n-m+1 .$$

Suppose A is an alphabet of size n and S is a sequence from $Q'(A,m)$ of minimum length (i.e. $|S|=M'(n,m)$). It is noted in (2.4) that $M'(n,1)=n$ so take $m\ge 2$. Segment S as $T\times U$ uhere the sequences T, U and the symbol x are chosen so that x does not appear in T but all the other symbols of A do. Clearly, $|T|\ge(n-1)$. Now note that al l permutations of subalphabets of A of size m which start with x are subsequences of xu. Hence all permutations of subalphabets of $A\backslash x$ of size $(m-1)$ are subsequences of U ($A\backslash x$ is A without x and $|A\backslash x|=(n-1)$). $|U|\ge M'(n-1,m-1)$, therefore, and so $M'(n,m)$ (which is simply $|S|$) is at least $(n-1)$ t 1 t $M'(n-1,m-1)$. This recurrence relation is readily solved to give the result.

2.7 $M(n)\ge(n.(n+1)/2)$.

Simple corol lary of 2.6' using $M(n)=M'(n,n)$.

2.8 $M'(n,m) \le (m.(n-1)+1)$

Given an alphabet, A, of size n , the following construction gives an element of $Q'(A,m)$ of length $m*(n-1)+1$:-
Generate m permutations of the 'alphabet Al, A2, A3, . . . Am such that Al(n)=A2(1), A2(n)=A2(1) etc. Now, B = Al A2(2:n) A3(2:n)...Am(2:n) is in $Q'(A,m)$ since if C is any permutation of any subalphabet of A of size. m, C(j) is either in the j-th component of B or IS the last symbol of the (j-1)th component (for j>1).

2.9 $M(n) \le (n.n-n+1)$

A simple corollary of 2.8.

2.10 $M'(n,3) = (3n-2)$ $(n \ge 3)$.

From 2.6 we get Il' $(n,3) \ge (3n-3)$.
From 2.8 we get $M'(n,3) \le (3n-2)$.
Suppose the lower value is obtained for an alphabet A ($|A|=n$) and S is a sequence of length 3n-3 which is in $Q'(n,3)$. Now no symbol can appear only once in S for then we would have $|S| \ge (2.M(n-1,2)+1) = (4n-5)$ which is a contradiction for $n \ge 3$. Hence there must be at least 3 symbols which occur just 2 times each for a total of 6 times. However $M(3)=7$ so there must be some permutation of these three symbols which is not a subsequence of S. This contradiction gives us the result.

2.11 Members of Q(1 2 3) of Length 7.

The fol lowing is an exhaustive list of minimum solutions for a 3 symbol alphabet. We consider, of course, only equivalence classes (with respect to the operator ≡).

1 2 3 1 2 1 3	1 2 3 1 2 3 1	1 2 3 1 3 2 1
1 2 3 2 1 2 3	1 2 3 2 1 3 2	
1 2 1 3 1 2 1	1 2 1 3 2 1 2	

2.12 $\forall S \in Q(A).\ \exists a \in A.\ \#(S,a) \ge |A|$.

Use induction on the alphabet size. The case $|A|=1$ is trivial so suppose the result holds for all alphabets of size less than n, $|A|=n$ and $S \in Q(A)$. Segment S as TxU where sequences T,U and symbol x are chosen so that x does not appear in T but every other symbol of A does. Use $A\backslash x$ to denote A minus symbol x, and ue get $U \in Q(A\backslash x)$. Now $|A\backslash x|$ = n-l and so we can find y such that $\#(U,y) \ge (n-1)$. Clearly $\#(S,y) \ge n$.

2.13 $\forall S \epsilon Q'(A,m) . Card(\{ a \mid a \epsilon A \wedge \#(S,a) \geq m \}) \geq (n-m+1)$

Let A be any alphabet, m be any integer such that $|A| \geq m$ and S be some member of $Q'(A,m)$. Select sequence B - a permutation of A such that the symbols of B are in order of decreasing frequency in S.
Now take sequence S' to be the sequence formed by deleting those symbols from S which are in $B(1:n-m)$. S' is a member of $Q(B(n-m+1:n))$ and so some symbol must appear at least m times in S' and hence in S.
Therefore, $\#(S,B(1)) \geq \#(S,B(2)) \geq \ldots\ldots \geq \#(S,B(n-m+1)) \geq m$ which gives the quoted result.

2.14 $M'(n,m) \geq m(n-m)+M(m)$

A corollary of 2.13 .

2.15 $M(4)=12$.

Take A to be the alphabet (sequence) 1 2 3 4 .

1 2 3 4 1 2 3 1 4 2 1 3 $\epsilon Q(A)$ and so $M(4) \leq 12$.

Suppose $S \epsilon Q(A)$ and $|S|<12$.
Compute the least integer j such that $S(1:j)$ contains each symbol of A. Note $j \geq 4$ and $S(j)$ is not in $S(1:j-1)$.
Considering permutations of A which start with $S(j)$, we get that $|S| \geq 3 + \#(S,S(j)) + M(3) = 18 + \#(S,S(j))$.
Using $|S|<12$ ue get $j=4$ and $\#(S,S(j))=1$.
Therefore, $S(4)$ appears only at position 4 of S. Now consider the permutations of A that end with $S(4)$ and get that $4 \geq M(3)$ which is a contradiction.

From this contradiction we see that $M(4) \geq 12$.

2.16 $\forall A . \forall x \epsilon A . \exists S \epsilon Q(A) . \#(S,x)=1$

Suppose we are given an alphabet A and x is **some** symbol of A. We take the subalphabet $A \backslash x$ and find some member T from $Q(A/x)$. Clearly $TxT \epsilon Q(A)$ and also $\#(TxT,x)=1$.
This is quite a useful result to keep in mind when pondering what properties members of $Q(A)$ might have.

3 M(5)=19.

Take A to be the alphabet (sequence) 1 2 3 4 5 .

i) 1 2 3 4 5 1 2 3 4 1 5 2 3 1 4 5 2 1 3 $\in$ Q(A)
so we have M(5) $\le$ 19 .

ii) Suppose S$\in$Q(A) and |S|<19 .
Break up S as T y U (where T and U are segments of S and y is a single symbol) such that Ty is the shortest initial segment of S which is in Q′(A,2) so |Ty| $\ge$ M′(5,2)=9.
Choose x in T such that xy is not a subsequence of T (this is possible otherwise S was not segmented as prescribed).

Considering members of P(A) starting with xy, get
-- |S| $\ge$ 3 t M(3) t #(U,x) t #(U,y) = 1 6 t #(U,x) t #(U,y).

Now, supposing x does not appear in U, consider subsequences of S that end with x and derive the contradiction
|s| $\ge$ M(4)+2+M(3)=21.
Conclude #(U,x) $\ge$ 1 (and similarly #(U,y) $\ge$ 1).

Reconciling inequalities, we get #(U,x)=1, #(U,y)=1, |T|=8, |U|=9 and |S|=18.

In U, x and y appear just once each and so one sequence of xy and yx , call it Z, is not a subsequence of U.
Consider, then, permutations of A of the form αZ and get
|T| $\ge$ M(3) t #(T,x) t #(T,y) $\ge$ 9 -- a contradiction!

We therefore conclude that M(5) $\ge$ 19.

iii) From i) and i i) deduce M(5)=19.

4 $M(6)=28$ and $M(7)=39$.

i) Take A to be the alphabet (sequence) 1 2 3 4 5 6 .

1 2 3 4 5 6 1 2 3 4 5 1 6 2 3 4 1 5 6 2 3 1 4 5 6 2 1 3
is in $Q(A)$ so we have $M(6) \leq 28$.

The proof of $M(6) \geq 28$ is given as Appendix 1 because it is. long and uninformative.

These two facts give the result $M(6)=28$.

ii) Take A to be the alphabet 1 2 3 4 5 6 7 .

1 2 3 4 5 6 7 1 2 3 4 5 6 1 7 2 3 4 s
1 6 7 2 3 4 1 5 6 7 2 3 1 4 5 6 7 2 1 3
is in $Q(A)$ so we have $M(7) \leq 39$.

$M(7) \geq 39$ (proved as appendix 2) and so we have $M(7)=39$.

5 Minimum Length Solutions for Alphabets of Size 4.

Let A be the alphabet **a b c d** .
We wish to enumerate the equivalence classes in $Q(A)$ of the minimum length (ie 12). Suppose $S \in Q(A)$ and $|S|=12$.

Lemma: $\forall p \in A.\ \#(S,p) \geq 2$
$p \in A \wedge \#(S,p)=0$ is absurd.
Suppose $p \in A \wedge \#(S,p)=1$ We have that S has the form UpV.
CPAF ap to get $|U| \geq M(3)=7$; CPAF pa to get $|V| \geq M(3)=7$.
We immediately have the contradiction $|S|=|UpV| \geq 15$.

Lemma: $\exists p.\ \#(S,p)=2$
Suppose not. In view of above lemma, $\forall p \in A. \#(S,p) \geq 3$ which is a violation of the result 2.12 (page 3).

Supposing $\#(S,p)=2$, choose T,U,V such that $S = TpUpV$.
CPAF pa to get $|UV| \geq 7$; CPAF ap to get $|TU| \geq 7$.
Now $|U| = |U|+(|S|-12) = (|U|+|T|+|U|+|V|+2)-12 \geq 4$.
Also $|T| = |S|-2-|U|-|V| \leq 3$ and similarly $|V| \leq 3$.

Suppose $|T|<3$. Thus $\exists x \in A.\ \neg(x \in T) \wedge \neg(x=p)$.
CPAF xpa to give $|V| \geq M(2)+\#(V,x)=3+\#(V,x)$. So $\#(V,x)=0$.
CPAF apx to give the contradiction $|T| \geq M(2)=3$.
Hence $|T|=3$ and similarly $|V|=3$ giving $|U|=4$.

Suppose $q \in A$ and $\neg(q=p) \wedge \#(T,q)=0$.
CPAF qpa to get $\#(V,q)=0$. Hence by a lemma above, $\#(U,q) \geq 2$.
CPAF qαp to get the contradiction $|U| \geq M(2)+\#(U,q) \geq 5$.
Hence $\forall q,\ q \in A \supset (q=p \vee \#(T,q)=\#(V,q)=1)$.

From this discussion we get that there are representatives of all the equivalence classes of the form
a b c d U d V where $|U|=4, |V|=3$, a d , $b \in V, c \in V$.

CPAF ad we get **abcU** is in $Q(a\ b\ c)$ and is of min. length.
Using result (2.11) we get 5 possibilities for U; namely:
(1) abac (2) abca (3) acba (4) babc (5) bacb.

Similarly UV is in $Q(a\ b\ c)$ and is of minimum length.
Performing a small amount of hand checking and using 2.11 again we get that there are exactly 9 equivalence classes:-

abcd abca dbac	abcd acba dbca	abcd bacb dabc
abcd abca dbca	abcd acba dcab	abcd bacb dacb
abcd abca dcba	abcd acba dcba	abcd bacb dcab

6. An $n^2 -2n +4$ Construction for Alphabet of size n.

Given an alphabet sequence, A, of length at least three, it is asserted that the following recipe gives a sequence in Q(A).

```
Set the sequence variable   B ← A(2:n);

Write(A):
DO (n-2) TIRES {Write(A(1));Write(B(1:n-2));
                B ← (B(n-1) B(1:n-2)); 3 ;
Write(A(1));  Write(B(1));
```

The total number of symbols written = n+ (n-2)*(1+n-2)+2
= n^2-2n+4 .

We now verify that the sequence produced is indeed in Q(A).

First note that the operation " B ← B(n-1)B(1:n-2)" simply rotates the sequence of n-l symbols in B.

Next note that the first symbol of A (we will call it a) is written exactly n times, Letting C be the result of the above construction, we segment C as follows:

C = aJaKaLa... aYaZab where the (n-1) sequences J,K,L,...Y,Z do not contain the symbol a.

For convenience we will use call J,K,L,.....Y,Z units and will refer to them as U[1] , U[2], . . . U[n-1].

Now J contains all symbols A(2:n) but K,L,...Y,Z each contain just n-2 of the symbols of A(2:n). However the symbol of A(2:n) that does not appear in some unit U[k] is both the last symbol of U[k-1] and follows the a that follows U[k] in C.

Let P be a permutation of A. We will show that P must be a subsequence of C.

Suppose a appears in the jth position of P. We first show that the string P(1:j) (simply a if j=1) can be matched to the the head of C aJaKaL...U[j-1]a . Trivially true if j=1.
If j>1 then P(1) is in J, clearly. Also if j>k>1 then P(k) can be matched to U[k] if it is in that unit or else the last symbol of U[k-1].
Similarly the n-j symbols of P(j+1:n) can be matched to U[j]aU[j+1]a...aU[n-1]ab . If j<k≤n then P[k] will either match some thing in U[k-1] or the symbol which follows the a which follows U[k-1].

7. A More General n^2-2n+4 Construction,

It is asserted that the follouing algorithm, regardless of which internal choices are made, also produces a member of Q(A) of length n^2-2n+4. The proof. of membership in Q(A) follows by the same method used in proving the validity of the simpler 'program'. It is also readily seen that the previous construction is a special case of this more general one.

SUBROUTINE SR1:
    Write the symbol [x];
    Write the symbol [y];

SUBROUT I NE SR2:
    SR1;
    Write in any order the [n-3] symbols of A which do not include
    [x] or [y] *or* [z].
    00 y←z AND set z to the last symbol written.

SUBROUTINE SR3: .
    DO SR2 k-21 TINES;
    SR1;

SUBROUTINE SR4:
    DO SR2 in-31 TIMES;
    SR1;
    Wr i te in any order the [n-2] symbols of A which are not [x],[y];
    Wri te the symbol [x];

MAIN ROUTINE:
    Write down the alphabet (A);
    DO EITHER {x←A(1); y ← any symbol of A(2:n-1);z←A(n);}
    OR { x ← A(2); y ← A(1); z ← A(n); };
    DO EITHER SR3 OR SR4;

SYMBOL COUNT.

If M symbols are written each time a certain routine is obeyed then we say that the SYMBOL COUNT for that routine is M.

Symbol Count for SR1 = 2 ;
Symbol Count for SR2 = n-l ;
Symbol Count for SR3 = $(n-2)*(n-1)+2 = n^2$ -3n +4;
Symbol Count for SR4 = $(n-3)*(n-1)+(n+1) = n^2$ -3n +4.
Hence Symbol Count for total algorithm = n^2 - 2n + 4 .

Note that no distinct sequences produced by this algorithm are equivalent since al l such begin uith a copy of the alphabet.

Note also that every sequence so produced ends with some permutation of the alphabet.

Given an alphabet A, the reversat of any sequence which is a member of Q(A) is also a member of Q(A). It should be noted that the the reverse of any sequence generated according to this construction is equivalent to some other sequence given by the construction.

8. Constructing Elements of $Q'(A,m)$.

Section 6 contained a simple construction for generating elements of $Q(A)$ (for given alphabet A of size $n>2$) which were of length n^2-2n+4. This algorithm is now modified to generate members of $Q'(A,m)$ (where $2<m\leq n$) of length $mn-2m+4$.

```
Set the sequence variable   B ← A(n-m+2:n);
Write(A);
DO m-2 TIRES Write(A(1:n-m+1));
             Write( B(1:m-2) );
             B ← B(m-1)B(1:m-2);
Write( A(1:n-m+1) );
Write( B(1) );
```

The total number of symbols written is easily seen to be

n t (m-2) (n-m+1 t m-2) t (n-m+1) t 1 = mn-2mt4 .

Just as this algorithm is a modification of the one in section 6, the proof of the correctness of the construction is an extension of the previous proof,

This construction gives an upper bound on $M'(n,m)$ for $n\geq m>2$ of $mn-2m+4$ and so using this knowledge, the proposition 2.14 and the various values of $M(4)$, $M(5)$, $M(6)$ & $M(7)$ we already know, we compute the neu results:-

$M'(n,4) = 4n-4$
$M'(n,5) = 5n-6$
$M'(n,6) = 6n-8$
$M'(n,7) = 7n-10$

9. Discussion,

The construction of section 7 gives many sequences of the desired length. It gives all nine equivalence classes of sequences in Q(a b c d) of length 12, 128 classes in Q(a b c d e) which may or may not be all of them, and 32,400 classes from Q(a b c d e f). It does NOT get all the sequences of Q(a b c d e f) since all the ones produced start with one copy of the alphabet however the following sequences from Q(a b c d e f):

abcdebfdcabedcfbadecbdfacebd

abcdeafdcbaedcfabdecafdbcead

(among others known) DO NOT! In fact, the second of these examples does not even end with a permutation of the alphabet.
An easy to derive lower bound on the number of classes is ((n-3)!)↑(n-1).

We now tabulate the known values of the functions M & M′ .

m	M(m)	m^2-2m+4	M′(n, m)
1	1	3	n
2	3	4	2n-1
3	7	7	3n-2
4	12	12	4n-4
5	19	19	5n-6
6	28	28	6n-8
7	39	33	7n-10

The fact that the actual values of M(n) exactly match the n^2-2n+4 figure for 2 < n ≤ 7 make the construction relatively important. It also suggests the obvious conjecture that M(n) is exactly n^2-2n+4 for all n>2. However, there is another competing conjecture which gives exact fit at n=1,2 as well as the other known values of M(n) but is more complicated:-

M(n) =	n^2	for n=1
	n^2-n+1	for 2≤n≤3
	n^2-2n+4	for 4≤n≤7
	$n^2-3n+11$	for 8≤n≤15
	,......	
	n^2-m*n+F(m)	for 2^m ≤n≤ 2.2^m-1

where F(0)=0 A F(n)=n+2*F(n-1).

Of course, knowing whether the value for M(8) is 51 or 52 would help by eliminating one of these postulates.

It is surprising that the best louer bound we have on $M(n)$ is $n^2/2$ since it would appear that it is of order n^2. This conjecture is readily stated formally as:-

$$\forall k.\ k<1 \supset \exists N.\ n>N \supset (M(n) > k*n^2)$$

It should be noted that just the mechanical checking of the membership of a sequence (over alphabet A) in Q(A) is quite time- consuming, A program is available in ALGOL but (although it includes some means for pruning the tree of permutations) takes a long time to check that all permutations of the alphabet are subsequences of the given sequence. The actual times on a PDP10 are 3, 17 and 60 seconds for alphabets of sizes 8, 9 & 10 respectively.

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1. Chvatal,V., Klarner,D.A., Knuth,D.E., "Selected Combinatorial Research Prob l ems", Report CS 292, Computer Science Department, Stanford University, June 1972.

APPENDIX 1. Proof of $M(6)\geq 28$.

Take A to be an alphabet of size 6 $(|A|=6)$.
Moreover, suppose $S\epsilon Q(A)$ and $|S|<28$.

Now choose sequences **T, V** and symbols x,y such that

a) Tx is the shortest head of S that is in $Q'(A,2)$;

b) yV is the shortest tail of S that is in $Q'(A,1)$;

Choose $w\epsilon T$ such that $w\neq x$ A $\neg(wx\subset T)$.

We have immediately that $|T|\geq 10$, $|V|\geq 5$ and from consideration of the elements of $P(A)$ of **the** forms $wx\alpha$ 6 βy get $|S|\geq|T|+1+M(4)$, $|S|\geq|V|+1+M(5)$, $|T|\leq 14$, $|V|\leq 7$, $|S|\geq 25$. Hence we can segment S as the sequence $TxUyV$ and note $10\leq|T|$ 114, $2\leq|U|$ 510, $5\leq|V|\leq 7$, $25\leq|S|\leq 27$.

Again CPAF wxa and get $|UyV|\geq M(4)+2=14$. Hence (using $|S|\leq 27$) $|T|\leq 12$ and (using $|V|\leq 7$) $|U|\geq 6$. Also CPAF ay again to deduce $|TxU|\geq M(5)+1=20$. Therefore, $|S|\geq 20+1+|V|\geq 26$ and (using $|T|\leq 12$) $|U|\geq 7$. Lastly (using $|S|\leq 27$ and $|TxU|\geq 20$), $|V|\leq 6$.

Suppose $\#(U,w)=0$. Since $|yV|\leq 7$ but contains all of A, there must be 5 symbols of yV which appear just once. Therefore we choose p,q such that p,q,x,w are distinct, $\neg(pq\subset yV)$ and p,q both appear **twice** in T. We can do this since only one symbol of Tx can appear only once. Now CPAF awpq to get $|T|\geq M(3)$ t $\#(T,w)$ t $\#(T,p)$ t $\#(T,q)\geq$ 1 2 , So $|T|=12$ and $\#(T,w)=1$. Segment S as $LwMxUyV$ noting that since $LwMx$ is in $P(A,2)$ and $\#(L,w)=0$, $|M|\geq 4$. This gives that $|L|\leq 7$ and $\#(MxU,w)=0$. $M(5,2)=9$ so we pick p,q such that $\neg(pq\subset L)$ and p,q,w distinct, Now CPAF pqwa to get $|yV|\geq M(3)+\#(yV,w)\geq 8$. **This contradiction gives** $\#(U,w)\geq 1$.

Again CPAF wxa and get $|UyV|\geq M(4)+\#(UyV,w)+\#(yV,x)\geq 15$. Use $|S|\leq 27$ to get $|T|\leq 11$ and use $|V|\leq 6$ to get $|U|\geq 8$.

. Now let $t\epsilon A$ be such that $\#(U,t)=0$. As above we choose p,q so that t,p,q are distinct, $\neg(pq\subset yV)$ and p,q both appear at least twice in T. CPAF atpq to deduce the contradiction

$|Tx|\geq M(3)$ t $\#(Tx,t)$ t $\#(Tx,p)$ t $\#(Tx,q)\geq$ 1 2 !!

Hence all symbols appear at least once in U.

Yet again CPAF wxa to get $|UyV|\geq M(4)+\#(UyV)+\#(UyV)\geq 16$. As before deduce $|T|\leq 10$ and $|U|\geq 9$. Al so CPAF ay to give $|TxU|\geq M(5)+\#(TxU,y)\geq 21$ and then $|S|=27$, $|V|=5$ We also have $|T|=10$, $|U|=10$ and $\forall t.\, t\epsilon A\supset t\epsilon U$.

The proof is concluded by deriving contradictions in the various possible cases of equality among w,x,y.

CASE **1.** $x=y$. and so **S** . $TxUxV$.
We know $\#(T,x)\geq 1$ and $\#(U,x)\geq 1$ **so CPAF** ax and get the **contradiction** 2 1 $=|TxU|\geq M(5)$ t $\#(TxU,x)\geq$ 2 2 .

CASE 2. $x \neq y$.

CASE 2a. $w \neq y$ i.e. w, x, y all distinct).

CPAF wxay to get $|U| \geq M(3)+\#(U,w)+\#(U,x)+\#(U,y) \geq 10$

Therefore $\#(U,w) = \#(U,x) = \#(U,y) = 1$.

Now this gives that one of wx or xw, call it Z, is such that $\neg(Z \subset U)$. CPAF $\alpha Z y$ and get $|T| \geq M(3)+\#(T,w)+\#(T,x)+\#(T,y)$

But $\#(T,w)+\#(T,y) \geq 3$ and so $|T| \geq 11$ -- contradiction!!

CASE 2b. $w=y$.

Find the first symbol of V which is not x ; call it z.

Note that since $yV \epsilon P(A) \wedge |yV|=|A|$, z appears just once in V.

CPAF yxaz to deduce $|U| \geq M(3)+\#(U,y)+\#(U,x)+\#(U,z) \geq 10$,

Immediately we see $\#(U,x)=\#(U,z)=1$ and so one of xz, zx (call it Z) is not a subsequence of U.

CPAF $\alpha Z y$ to get $|T| \geq M(3)+\#(T,x)+\#(T,y)+\#(T,z)$.

'Use $\#(T,y)+\#(T,z) \geq 3$ for the contradiction $|T| \geq 11$.

APPENDIX 2. Proof of $M(7)\geq 39$.

Take A to be an alphabet of size 7 $(|A|=7)$.
Moreover, suppose $S\in Q(A)$ and $|S|<39$.

Choose sequences T,U,W and symbols a,b,c such that
a) Ta is the shortest head of S that is in $Q'(A,1)$
b) cW is the shortest tail of S that is in $Q'(A,1)$
c) $TaUb$ is the shortest head of S that is in $Q'(A,2)$

We segment S as $TaUbVcW$ and readily prove:
$6\leq|T|\leq 8,\ 5\leq|U|\leq 9,\ 8\leq|V|\leq 18,\ 6\leq|W|\leq 8,\ 36\leq|S|\leq 38;$
as well as $|T|+|U|\leq 15$.

Suppose for some p in A, $\#(V,p)=0$.
If p is the symbol b, $M'(6,3)+\#(TaUb,p)\geq 18>|TaUb|$ so we can choose q,r,s such that $distinct(p,q,r,s)\wedge\neg(qrs\subset TaUb)$ so that $\neg(qrsp\subset TaUbV)$. CPAF qrspa we get a contradiction $|cv|\geq 4+M(3)$.
Otherwise p,b are distinct and $M'(6,3)+\#(TaU)\geq 17\geq|TaU|$ so we rechoose q,r,s such that $distinct(p,q,r,s)\wedge\neg(qrs\subset TaU)$ which means $\neg(qrsp\subset TaUbV)$. As before get a contradiction.
Lemma 1: $\forall x\in A.\ \#(V,x)\geq 1$ follows from these contradictions.

Suppose $p\in A\wedge distinct(a,p)$. We know $\#(T,p)\geq 1$ and $\#(Ub,p)\geq 1$ and $\#(V,p)\geq 1$ and $\#(cW,p)\geq 1$ so conclude $\#(S,p)\geq 4$. Also we have $\#(V,a)\geq 1$ and $\#(cW,a)\geq 1$ so that $\#(S,a)\geq 3$.
We sharpen our inequalities now. CPAF $a\alpha$ to get $|T|\leq 7, |S|\geq 37$; CPAF aba to get $|T|+|U|\leq 13$; CPAF ab to get $|W|\leq 7$. Hence
$6\leq|T|\leq 7,\ 5\leq|U|\leq 7,\ 13\leq|V|\leq 18,\ 6\leq|W|\leq 7,\ 37\leq|S|\leq 38.$

Suppose, in fact, $\#(S,a)=3$.
We re-segment S as $TaJaKaL$ where $\#(TJKL,a)=0$ and $L\subset W$.
There is at most one repeated symbol in T since $|Ta|\leq|A|+1$.
Let z denote this symbol if it exists else any symbol of T.
Choose p,q such that $distinct(p,q,a,z)\wedge\neg(pq\subset T)$.
CPAF pqzaa to deduce that some subsequence G of KaL belongs to $Q(A1)$ where A1 is obtained from A by deleting p,q,a,z.
$|G|\geq M(3)=7$ so some symbol of G appears at least 3 times.
So we choose y to be such a symbol and note
$distinct(a,y)\wedge\#(T,y)=1\wedge\#(KaL,y)\geq 3.$
Now one of py and yp (call it Z) is not a subsequence of T.
CPAF Zzaa to show we can choose x with the properties
$distinct(x,y,a)\wedge\#(T,x)=1\wedge\#(KaL)\geq 3.$

Now, one of the sequences xy and yx is not a subsequence of T (call it Y) and CFAP Yaa to get

$|KaL| \geq M(4) + \#(KaL,a) + \#(KaL,x) + \#(KaL,y) \geq 19$.

By symmetry $|TaJ| \geq 19$ to give the contradiction $|S| \geq 19+19+1$.

Lemma 2: $\forall x \epsilon A.\ \#(S,x) \geq 4$ is immediate.

Again CPAF $a\alpha$ to get $|T|=6$, $|S|=38$, $\#(S,a)=4$;
Also CPAF ac to deriv e $|W|=6$, $|U|+|V|=23$, $\#(S,c)=4$.
Then CPAF aba to get $|VcW| \geq M(5)+\#(VcW,a)+\#(VcW,b) \geq 23$
which leads to $16 \leq |V| \leq 18$ anti $5 \leq |U| \leq 7$.

Suppose that p,q are such that $\neg(pq \subset V)$. We have that $\#(TaUb,p)+\#(TaUb,q) \geq 3$. Now $|TaUb| \leq 15$ and so

$|TaUb| < M'(5,3)+\#(TaUb,p)+\#(TaUb,q)$. Hence we

choose j,k,l such that $distinct(j,k,l,p,q) \wedge \neg(jkl \subset TaUb)$.
CPAF jklpqa so $|cW| \geq M(2)+5=8 > |cW|$ -- a contradiction!
Thus $\forall p \epsilon A.\ \forall q \epsilon A.\ \#(V,p)+\#(V,q) \geq 3$.
In particular, letting zbe the first symbol of cW which is not one of a,b , $\#(V,a)+\#(V,b)+\#(V,z) \geq 5$.
CPAF $ac\alpha\bar{z}$ to get $|V| \geq M(4)+\#(V,a)+\#(V,b)+\#(V,z) \geq 17$
Thus we have new bounds for U,V:- $5 \leq |U| \leq 6$, $17 \leq |V| \leq 18$.

We now choose sequence H and symbol d such that

dHcW is the shortest tail of S in Q(A).

By symmetry with the results for U we have that $5 \leq |H| \leq 6$ and so we re-segment S as TaUbGdHcW where
$|T|=6$, $5 \leq |U| \leq 6$, $10 \leq |G| \leq 12$, $5 \leq |H| \leq 6$, $|W|=6$, $|S|=38$, $\#(S,a)=4$, $\#(S,c)=4$.

Suppose x is such that $x \neq a \wedge x \neq c \wedge \neg(e \epsilon G)$.
If $x \neq b$ then CPAF $abe\alpha$ to get

$|dHcW| \geq M(4)+(\#(dHcW,a)+\#(dHcW,b))+\#(dHcW,e) \geq 12+3+2$

- a contradiction.

If xzd then CPAF aedc to get

$|TaUb| \geq M(4)+(\#(TaUb,c)+\#(TaUb,d))+\#(TaUb,e) \geq 12+3+2$

- also a contradiction.

The remaining case is $x=b=d$. Lemma 1 (with $\#(S,c)=4$) gives that $\#(TaUb,c) \leq 2$ and since there is at most one symbol in TaUb appearing 3 times, we choose p,q (not c or b) so that $\#(TaUb,p) \leq 2$ and $\#(TaUb,q) \leq 2$. Since $M(3)=7$ there is some permutation Z of c,p,q that is not a subsequence of TaUb. CPAF Zba to get
$|HcW| \geq M(3)+\#(HcW,b)+\#(HcW,c)+\#(HcW,p)+\#(HcW,q) \geq 7+1+2+2+2 = 14$.
- a contradiction.

From these 3 contradictions we get $(x \epsilon A \wedge x \neq a \wedge x \neq c) \supset \#(G,x) \geq 1$.
Now suppose $\neg(a \epsilon G)$. Choose p,q,r so that $distinct(a,p,q,r)$ and $\neg(pqr \subset dHcW)$. CPAF aapqr. Clearly $a \epsilon U$ [else $|T| \geq M(4)$] and so $\#(TaUb,a) \geq 2$. Hence

$|TaUb| \geq M(3)+\#(TaUb,a)+....+\#(TaUb,r) \geq 7+2+2+2+2 = 15$

From this contradiction we get $\#(G,a) \geq 1$ and by symmetry $\#(G,c) \geq 1$.

Lemma 3: $\forall x \epsilon A.\ \#(G,x) \geq 1$ follows.

Suppose $x \in A \wedge x \neq a \wedge x \neq c$. $\#(T,x)=\#(W,x)=1$, $\#(Ub,x)\geq 1$, $\#(dH,x)\geq 1$ and $\#(G,x)\geq 1$ to yield

Lemma 4: $\forall x \in A. (x \neq a \wedge x \neq c) \supset \#(S, x)\geq 5$.

Suppose distinct(a,b,c).
We first choose z to be the first symbol of W which is not a,b.
$b \neq a \wedge b \neq c$ so we have $b \in G, b \in dH$ giving $\#(GdH,b)\geq 2$.
$z \neq a \wedge z \neq c$ so we have $z \in G, z \in dH$ giving $\#(GdH,z)\geq 2$.
Also $a \neq c$ so $a \in dH$ and we have $a \in G$ giving $\#(GdH,a)\geq 2$.
CPAF abaz to derive $|GdH| \geq M(4)+\#(GdH,a)+\#(GdH,b)+\#(GdH,z) \geq 18$.
We get from this that $|U|=5$ and also $\#(GdH,b) = 2 = \#(GdH,z)$.
This then gives that $\#(S,z)=5$ and $\#(S,b)=5$.

Let p,q,r be the 3 symbols of the A which are not a,b,c,z.
$\#(S,a) + \#(S,b) + \#(S,c) + \#(S,z) = 4+4+5+5 = 18$
so $\#(S,p) + \#(S,q) + \#(S,r) = 28$.
Since no symbol appears twice in TaUb, can choose a permutation Z of pqr so that $\neg(Z \subset TaUb)$.
CPAF Zα to get $25=|GdHcW|\geq M(4)+(20-6)=26$ - a contradiction,
Similarly 'distinct(a,d,c)' gives a contradiction.

Lemma 5: $\neg$distinct(a,b,c) $\wedge$ $\neg$distinct(a,d,c).

In view of lemma 5, two important cases are $a=c$ and $\neg(a=c)$.

CASE 1. $a=c$.
Suppose first that $a \in U$. Clearly $|U|=6$ and $|TaUb|=14$.
Letting z be the first symbol of W not a,b CPAF abaz to get $|GdH| \geq 12+\#(GdH,a)+\#(GdH,b)+\#(GdH,z) \geq 17$.
But $|GdH|=17$ so we see $\#(GdH,b) = 2 = \#(GdH,z)$.
Thus $\#(S,a)+\#(S,b)+\#(S,z)=14$.
Now choose p,q,r,s such that pqrsabz is a permutation of A and $\#(S,p)\geq\#(S,q)\geq\#(S,r)\geq\#(S,s)$. Now since some symbol appears at least 7 times in S, $\#(S,p)\geq 7$ and $\#(S,q)+\#(S,r)+\#(S,s)\leq 17$.
Hence $\#(S,s)\leq 5$ and so $\#(S,p)+\#(S,q)+\#(S,r)\geq 19$.
Now each of p,q,r appears exactly twice in TaUb and so

i) $\#(GdHaW,p)+\#(GdHaW,q)+\#(GdHaW,r) \geq 13$

ii) since $M(3)=7$ there is a permutation of pqr (call it Z) such that $\neg(Z \subset TaUb)$.

CPAF Za to get $24=|GdHaW|\geq M(4)+13 = 25$.
This contradiction gives us $\#(U,a)=0$.

Again letting z be the first symbol of W not a,b we have $\#(GdH,a)\geq 2$, $\#(GdH,b)\geq 2$, $\#(Gdh,z)\geq 2$ so CPAF abaz to deduce $|GdH|\geq 18$ and hence $|U|=5$ and $\#(S,b)=\#(S,z)=5$
Similarly, $\#(S,d)=5$ and $|H|=5$.

$|G|=12$ and $\#(G,a) = \#(G,b)=2$ so the other 5 symbols appear a total of 8 times in G. Hence choose p,q so that $\neg(pq \subset G)$ and distinct(a,b,p,q). $\neg(abpq \subset TaUbG)$ so CPAF abpqa to derive a contradiction $|dHaW|\geq 7 + 3*2 + 1 = 14$, CASE 2. $\neg(a=c)$.

We have $a\neq b$ and $c\neq d$ so Lemma 5 gives both $b=c$ and $d=c$,
Hence S looks like TaUbGaHbW with $|T|=6, 5\leq|U|\leq 6, 10\leq|G|\leq 12$, $5\leq|H|\leq 6$, $|W|=6$, $\#(G,a)=\#(G,b)=1$, $\#(T,b)=\#(W,a)=1$.
Clearly $\#(TUH,a)=0=\#(UHW,b)$.

We can write the alphabet in order of decreasing frequency in S as pqrstab where all except a,b occur at least 5 times and $\#(S,p)\geq 7$. Hence, as p,q,r,s,t appear a total of 30 times $\#(S,t)=5$ and $\#(S,s)\leq 6$ and $\#(S,p)+\#(S,q)+\#(S,r)\geq 19$.

CASE 2a: $|U|=5$.

Some permutation, Z, of pqr will not be a subsequence of TaUb so CPAF Za to get $|GaHbW|\geq 12+19-6=25$.
This gives us that $\#(S,p)+\#(S,q)+\#(S,r)=19$ and $\#(S,s)=6$.
We then deduce $\#(S,p)=7, \#(S,q)=\#(S,r)=6$.

Now if z denotes the last symbol of T then CPAF za to get $32=|aUbGaHbW|\geq M(6)+\#(S,z)-1$ or $\#(S,z)\leq 5$
But $z\neq a$ so $\#(S,z)\geq 5$ so we deduce t-t.
Similarly the first symbol of W is t.

Recall that $\neg(Z\subset TaUb)$, $\#(G,a)=\#(G,b)=1$ and note $\#(G,t)=1$.
CPAF Zabα to deduce that ab $\subset$ G.
CPAF Ztba to deduce that tb $\subset$ G.
Similarly deduce that at $\subset$ G.
i.e. a precedes t precedes b (in G).

Suppose t is not the last symbol of U. We find y,z such that $\neg(yzt\subset TaUb)$ and so $\neg(yztab\subset TaUbGaH)$. CPAF yztab for the contradiction by which we can conclude $U(5)=t$.

We have that S has the form T′taU′ftbGaHbtW′ where T't-T, U′ft=U and tW′=W (this defines T', U', f, W′).
Clearly $f\neq a$, $f\neq b$, $f\neq t$ and so $\#(S,f)\geq 6$.
Now $\neg(tf\subset TaUb)$ so CPAF tfaab to get $|G|\geq 7+3+\#(G,f)$.
Suppose $\#(G,f)=1$. From $\#(S,f)\geq 6$ deduce $\#(H,f)=2$.
Now one of tf, ft is not in G - call it Z.
CPAF abZα to get $|aHbW|\geq 7+1+2+2+3=15$ - a contradiction.
Hence we have $\#(G,f)=2$ and $|G|=12$ so $|H|=5$.

Now let the last symbol of T' be g and suppose $b\neq g$.
$\neg(gb\subset TaU)$ and $\neg(ta\subset G)$ so $\neg(gbta\subset TaUbG)$.
CPAF gbtaa to get a contradiction,
Hence the last symbol of T' is b.

Now $\neg(bf\subset T'taU')$ but we have $\neg(ta\subset bG)$ so $\neg(bfta\subset TaUbG)$.
CPAF bftaa to get $12=|HbW|\geq 7+1+1+2+2=13$,
This last contradiction dispenses with CASE 2a.

CASE 2b: $|H|=5$.

The elimination of this case is similar to CASE 2a.

CASE 2c: $|U|\neq 5 \wedge |H|\neq 5$.

We have so far that $S = TaUbGaHbW$ with $|T|=|U|=|H|=|W|=6$
$|G|=10$, $\#(G,a)=\#(G,b)=1$, $\#(TUH,a) = \#(UHW,b) = 0$.

Suppose first that $\#(S,s)=5$.
Without loss of generality suppose s precedes t in G.
$\neg(abts \subset TaUbGa)$. Moreover if any p, q or r precedes s in H
then CPAF abtsa to get $|HbW|>7+1+1+4=13$ - a contradiction.
Hence only t may precede s in H.
Similarly only s may follow t in U.
Now CPAF atasb to get $|G|\geq M(7)+\#(G,a)+\#(G,b)+\#(G,s)+\#(G,t)=11$.
The contradiction serves to give us $\#(S,s)\neq 5$.
Hence $\#(S,s)=6$ and $\#(S,p)=7$, $\#(S,q)=\#(S,r)=6$.

Letting--.x be the duplicated symbol in U and y the dupl icatetd
symbol in H, $\#(U,x)=2$, $\#(H,y)=2$.
If x-y then $\#(S,x)\geq 7$ so x-p and thus $\#(G,x)=1$.
One of yt, ty (cal l it Z) is not a subsequence of G.
CPAF $abZ\alpha$ to get $|HbW|\geq 7+1+1+2+3=14$ - contradiction.
Else if $y\neq p$ then $\#(S,y)=6$ (note $y\neq a$, yrb, $y\neq t$) and $\#(G,y)=1$
One of yt, ty (cal l it Z) is not a subsequence of G.
CPAF $abZ\alpha$ to get $|HbW|\geq 7+1+1+2+3=14$ - contradiction.
Else $x\neq y \wedge$ y-p so $x\neq p$ and $\#(S,x)=6$.
One of xt, tx (cal l it Z) is not a subsequence of G.
CPAF αZab to get $|TaU|\geq 7+1+1+2+3=14$ - contradiction.

This trio of contradictions completely eliminates CASE 2c.

CASES 2a, 2b, 2c all provided contradictions as did CASE 1
so the assumption that $|S|<39$ is proved impossible.

Q.E.D.